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# Toward Reliable Multi-Level Operation in RRAM Arrays: Improving Post-Algorithm Stability and Assessing Endurance/Data Retention

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**ABSTRACT** Achieving a reliable multi-level operation in resistive random access memory (RRAM) arrays is currently a challenging task due to several threats like the post-algorithm instability occurring after the levels placement, the limited endurance, and the poor data retention capabilities at high temperature. In this paper, we introduced a multi-level variation of the state-of-the-art incremental step pulse with verify algorithm (M-ISPVA) to improve the stability of the low resistive state levels. This algorithm introduces for the first time the proper combination of current compliance control and program/verify paradigms. The validation of the algorithm for forming and set operations has been performed on 4-kbit RRAM arrays. In addition, we assessed the endurance and the high temperature multi-level retention capabilities after the algorithm application proving a 1 k switching cycles stability and a ten years retention target with temperatures below 100  $^{\circ}$ C.

**INDEX TERMS** RRAM, arrays, algorithm instabilities, multi-level, data retention, accelerated test.

#### I. INTRODUCTION

The resistive random access memory (RRAM) technology based on metal-oxide dielectrics is one of the most promising candidates for the next generation of non-volatile memory (NVM) applications due to their potential for high-density integration, high-speed switching operations, low-power consumption, and full compatibility with the CMOS technology [1]. The functionality of this technology is based on the resistive switching (RS) effect, which is attributed to the creation and disruption of nanometer-scale conductive filaments (CFs) in the insulator layer, consisting of oxygen vacancies  $(V_0)$  [2]. The creation process moves the RRAM device in a low resistive state (LRS), whereas the disruption process brings it in a high resistive state (HRS), which are referred as set and reset operations, respectively [3]. To activate the switching behavior, a preliminary soft breakdown in the dielectric material, referred as forming operation,

is required in most RRAM technologies [4], [5]. Such operation plays a crucial role in the device performance [6], [7].

Because of the stochastic nature of the CF formation and rupture, resistive states and their related switching parameters usually vary in a large scale (from device-to-device and from cycle-to-cycle), which compromises the RRAM reliability [8], [9]. Several methods have been proposed to improve the RS stability in RRAM devices, such as: active electrode modification [10], [11], metal-oxide interfaces engineering [12], doping of the switching layer [13], [14] or by using optimized programming algorithms [15], [16]. One of our prior works demonstrated a drastic reduction of the device-to-device variability by combining Al-doped HfO<sub>2</sub> as switching layer with an optimized programming algorithm [17].

That reduction paved the way for the introduction of the multi-bits per cell paradigm that is usually achieved

TABLE 1. Comparison with different multi-level programming algorithms in RRAM technology presented in literature and their associated reliability assessment tests.

Ref	Algorithm	N° of	Test	Reliability
	foundation		rest	avaluations
	Toundation	levels	venicle	evaluations
[18]	ISPVA-like	4 - 8	Single cell	Endurance
				Read disturb.
				Low temp. retention
[19]	Compliance	8	Single cell	Read disturb.
	control			Low temp. retention
[20]	Dispersion	4	Single cell	-
	aware PVA			
[21]	ISPVA-like	16	Single cell	Low temp. retention
[22]	ISPVA-like	92	Single cell	-
[23]	HRS control	4	10 cells	Endurance
[24]	Compliance	4 - 8	$2\times$	Endurance
	control with		4-kbit	
	smart target		arrays	
This	M-ISPVA (LRSs)	4	4-kbit	Post-algorithm
work	ISPVA (HRS)		arrays	instability
				Endurance
				Low and high
				temp. retention

by tailoring the CF properties through multi-level programming algorithms [18]-[24]. Table 1 summarizes the most common approaches found in the literature. However, there are a number of intrinsic concerns that could limit the reliable use of a multi-bit approach. Among them, there is a post-algorithm relaxation of the LRS and HRS distributions seen in several RRAM technologies [25], a difficult control of the multi-level endurance and a relatively poor high temperature data retention. Those issues have a consequence on the realization of high density products, confining the usage scope of RRAM technology in few scenarios. Further, the use of RRAM devices in neuromorphic applications, where the mimicking of biological synapses requires an analog-like behavior and therefore a stable multi-level capability [26], would be threatened a priori.

In this work, we demonstrate a well-controlled and reliable multi-level operation in 1-transistor-1-resistor (1T1R) RRAM architectures based on Al-doped HfO2 and integrated in 4-kbit arrays through the introduction of a new degree of freedom in the state-of-the-art incremental step pulse with verify algorithm (ISPVA), namely, the compliance current control, along with the target read-out current  $(I_{trg})$ . To the best of our knowledge, this is the first attempt to combine these two parameters in order to improve the reliability of multi-level RRAM technology, which will lead the way towards the multi-bit operation as well as the implementation of artificial synaptic devices [27]. The electrical characterization of the new algorithm, hereafter called multilevel-ISPVA (M-ISPVA), on RRAM array test vehicles will allow the investigation of the post-algorithm instability, memory endurance, and low/high temperature data retention capabilities. Finally, retention accelerated tests are performed to project the lifetime reliability up to 10 years as a further assessment for our developed algorithm on arrays.



**FIGURE 1.** Schematic illustration of the ISPVA (HRS) and of the M-ISPVA (LRSs) approaches (left), and I-V characteristics of the transistor setting the compliance current (right).

## II. EXPERIMENTAL

The test vehicles used for the implementation of the multilevel approach are 4-kbit memory arrays organized in 64 pages, each consisting of 64 1T1R cells. The architecture of the chip is described in more detail in [28]. The 1T1R devices are constituted by a NMOS transistor, manufactured in 0.25  $\mu$ m CMOS technology, whose drain is connected in series to a variable resistor integrated on the metal line 2 of the CMOS process. Such resistor is a metal-insulator-metal (MIM) structure consisting of a TiN/Al:HfO<sub>2</sub>/Ti/TiN stack with 150 nm TiN top and bottom electrode layers deposited by magnetron sputtering, a 7 nm Ti layer (under the TiN top electrode) and a 6 nm Al:HfO<sub>2</sub> layer based on  $\sim 10\%$  Aldoped HfO<sub>2</sub> (Hf<sub>1-x</sub>Al<sub>x</sub>O<sub>y</sub>) grown by atomic layer deposition (ALD) at 300 °C. After patterning the MIM structure with an area of about 0.4  $\mu$ m<sup>2</sup>, an additional thin Si<sub>3</sub>N<sub>4</sub> layer was deposited to protect the RRAM cell.

The key of a good multi-level approach in RRAM devices is in the accurate control of the multiple conductive states. The strategy followed in this work is to define one HRS and three LRSs, which leads to four levels cells. The ISPVA has shown to be an effective way to reduce the device-to-device variability by establishing  $I_{trg}$  values: a top value for HRS and a bottom value for LRS [17]. As schematically depicted in Fig. 1, during the ISPVA programming a sequence of increasing voltage amplitude pulses (P) of 10  $\mu$ s duration are applied either on the drain side of the transistor connected to the MIM resistor during forming and set operations or on the source terminal of the transistor during reset operations. After every pulse (P) a read-verify operation (V) is performed by applying a read-out pulsed voltage of 0.2 V amplitude and 10  $\mu$ s duration on the drain side of the transistor to check whether the  $I_{trg}$  target is achieved. In positive case the programming operation is stopped and the read-out current measured is saved as the read-out current value just after the transition. If the  $I_{trg}$  target is not achieved, the programming operation is stopped when a maximum voltage amplitude value is achieved: 3.5 V for reset and set operations and 5 V for the forming operation.

However, the definition of  $I_{trg}$  as a lower limit for LRS in the ISPVA approach is found deficient in ensuring three discrete levels as desired in our multi-level paradigm. As already claimed by Kim et al. [21], the abrupt nature of forming and set processes could easily program the RRAM devices into a conductive state outside of the desired range. Therefore, an upper target must be defined along with  $I_{tro}$ by using the current compliance control provided by the transistor in the 1T1R architecture, thus defining the M-ISPVA algorithm for LRSs. In the test vehicles considered in this work, the compliance current is effectively modulated by the gate voltage  $(V_g)$  applied on the transistor, whose I-V characteristics are shown in Fig. 1. Three  $\langle I_{trg}, V_g \rangle$  pairs were defined to evaluate the multi-level capabilities during forming:  $<20, 1.2>, <30, 1.4>, and <40, 1.6> <\mu A, V>.$ In the following we will refer to the LRS levels generated by those pairs as LRS1, LRS2, and LRS3, respectively. It is worth to point out that in order to ensure a good HRS, an  $I_{trg}$  value equal to 5  $\mu A$  was considered along with a  $V_g = 2.7$  V exploited in the ISPVA algorithm for the reset operation to minimize the series resistance of the transistor.

## **III. RESULTS AND DISCUSSION**

### A. FORMING OPERATION

In order to activate the RS behavior, the forming operation was performed by using each LRS  $\langle I_{trg}, V_g \rangle$  pair on different batches of pages of the array. The amplitude of voltage pulses applied on the transistor's drain during the M-ISPVA was swept in the range of 2-5 V with a step of 0.01 V. The three post-algorithm cumulative distribution functions (CDFs) of the read-out currents just after the forming transition and at the end of the M-ISPVA are shown in Fig. 2. An architecture constraint of the test setup imposes read-verify operations (V) to be continuously performed on the array until the M-ISPVA finishes (even on cells already formed). The read-out current measured during the last read-verify operation of the whole sequence is saved as the read-out current value at the end of the operation. The time interval between the switching transition of the RRAM cells and the end of the algorithm is averagely 3 minutes. This leads to the observation of post-algorithm instabilities yielding to perturbations of the conductive state of the RRAM devices under test [25]. The most likely reason for the occurrence of such instabilities is the further electrical stress applied by the additional read-verify operations. This is appreciable in the figure showing a lower tail of the CDFs with about 13% of the cells redistributing their read-out current below the  $I_{trg}$  and about 3% of the cells crossing the adjacent LRS level. This result is in accordance with [29]. Fig. 3 illustrates the device-to-device correlation between the current values measured just after the forming transition and at the end of the M-ISPVA. The lack of correlation between the current values shows that the tailing and the redistribution are stochastic processes whose detrimental effects on the multi-level operation should be prevented.



**FIGURE 2.** CDFs of the read-out currents just after the forming transition (full symbols) and at the end of the forming operation (empty symbols) for the three LRS levels : <20, 1.2> (in blue), <30, 1.4> (in green), and <40, 1.6> (in red) < $\mu$ A, V>.



**FIGURE 3.** Device-to-device correlation of the read-out currents measured just after the forming transition and at the end of the M-ISPVA for the three LRS levels: <20, 1.2> (in blue), <30, 1.4> (in green), and <40, 1.6> (in red) < $\mu$ A, V>.

The definition of suitable values for  $\langle I_{trg}, V_g \rangle$  pairs in the M-ISPVA is fundamental in providing multi-level capability while preventing post-algorithm instabilities, therefore outperforming ISPVA [17] and compliance-based algorithms [19] in this context. Multiple combinations of  $\langle I_{trg},$  $V_g$  > were tested during the forming operation besides those exploited for former tests. Three  $I_{trg}$  values were considered, namely 20, 30, and 40  $\mu A$  in combination with four Vg values: 1.0, 1.2, 1.4, and 1.6 V. The results of these measurements are summarized in Fig. 4. On the one hand, if the  $V_g$  value is too small, the RRAM cells can not achieve the target value I<sub>trg</sub> performing the forming transition. In this specific situation the current value saved as the read-out current just after the transition is the value measured at the end of the forming operation. This is the reason why both CDFs appear partially or even totally overlapped. On the other hand, the dispersion of the CDF increases beyond control



**FIGURE 4.** CDFs of the read-out currents just after the forming transition (solid symbols) and at the end of the forming operation (open symbols) for the twelve  $< l_{trq}$ ,  $V_g >$  pairs.



**FIGURE 5.** Evolution of the read-out forming current values during the M-ISPVA by using the <30, 1.4>  $\langle \mu A, V \rangle$  pair on 128 RRAM devices (in grey). The figure highlights the evolution on a cell reaching the forming  $l_{trg}$  with one step (in dark green), and a cell requiring more than one step to reach the forming  $l_{trg}$  (in bright green).

with large  $V_g$  values, which makes difficult to fit several conductive levels within the available read-out current range (from 0 to 50  $\mu$ A).

To explain this behavior, the evolution of the read-out currents measured during the M-ISPVA forming operation was carefully investigated. For instance, we will consider the batch of 128 RRAM cells formed by using the <30,  $1.4 > <\mu A$ , V> pair (see Fig. 5). In the pristine state (in which the read-out current is essentially equal to 0  $\mu A$ ) almost the whole voltage applied by the M-ISPVA on the RRAM device drops on the MIM cell. When a soft breakdown happens in the device's dielectric, a sudden increase of the read-out current takes place. According to Ielmini [30], the amplitude of such increase is limited by the voltage divider formed between the transistor resistance (controlled by V<sub>g</sub>) and the MIM cell resistance. However, a successful forming transition crossing I<sub>trg</sub> is not granted for all devices,



**FIGURE 6.** CDFs of the read-out currents at the end of the M-ISPVA during the forming operation for all  $<|_{trg}$ ,  $V_g >$  pairs.

since some of them requires additional voltage steps to converge to the target current. In addition, even if the forming transition is performed, the post-algorithm instabilities are appreciable. Such instabilities are totally independent on the  $I_{trg}$  value defined and can be limited only by the current compliance control imposed by  $V_g$ , as shown in Fig. 6. This is an added value provided by the M-ISPVA compared to other algorithms.

# B. RESET AND SET OPERATIONS

After forming the three LRS levels, the feasibility of switching the RRAM devices from the LRS1, LRS2, and LRS3 to the HRS and viceversa was tested by performing reset and set operations with ISPVA and M-ISPVA, respectively. In both algorithms, the amplitude of voltage pulses applied on the source (reset) and drain (set) was swept in the range of 0.2-3.5 V with a voltage step of 0.1 V. As shown in Fig. 7, the CDFs of the read-out currents after the reset operation are completely located below the  $I_{tre}$  defined for the HRS (5  $\mu$ A) in the ISPVA algorithm. After reset, the three LRSs were obtained after the set operation with M-ISPVA using the same  $\langle I_{trg}, V_g \rangle$  pairs defined for the forming operation. About 5% of the devices for each LRS level cross the  $I_{trg}$  value of the adjacent level at the end of the algorithm. The CDF corresponding to the LRS3 displays a tail at high current values, however this does not interfere with the multi-level definition. Such tail can be attributed to a proximity with the upper limit of the read-out current range (50  $\mu$ A) where the switching behavior of our devices starts to be unstable. Unlike in the forming operation, in set operation the redistribution of the current values during the time interval between the switching transition and the end of the algorithm (30 seconds in average) is almost suppressed, leading to a more reliable multi-level operation (see Fig. 8). An explanation of this phenomenon can be found in [31]. At the beginning of the first reset operation, the oxygen vacancies placed at the interface between the



**FIGURE 7.** CDFs of the read-out currents after the reset operation with ISPVA (HRS) and after the set operation with M-ISPVA (LRSs). LRSs data are shown after the set transition (solid symbols) and at the end of the algorithm (open symbols) : <20, 1.2> (in blue), <30, 1.4> (in green), and <40, 1.6> (in red) < $\mu$ A, V>.



**FIGURE 8.** Device-to-device correlation of the read-out currents measured just after the set transition and at the end of the M-ISPVA for the three LRS levels: <20, 1.2> (in blue), <30, 1.4> (in green), and <40, 1.6> (in red) < $\mu$ A, V>.

Al:HfO<sub>2</sub> insulator and the bottom electrode drift towards the CF, thus strengthening the CF tip in the first voltage steps of the ISPVA algorithm applied during reset operation. Such increase, independent on the starting LRS level, vanishes when some of the oxygen vacancies drift and open a gap in the CF. The subsequent set operation will fill the gap previously opened only by a few oxygen vacancies. This process stabilizes the CF strongly reducing the post-algorithm instabilities for further operations.

As for forming operation, it is important to define appropriate  $\langle I_{trg}, V_g \rangle$  pairs to establish multiple LRSs in RRAM devices without endangering their reliability with unoptimized M-ISPVA conditions. Fig. 9 shows the CDFs of the read-out currents just after the set transition and at the end of the M-ISPVA for the same twelve  $\langle I_{trg}, V_g \rangle$  pairs considered in the forming analysis. Similarly, if the  $V_g$  value



**FIGURE 9.** CDFs of the read-out currents just after the set transition (solid symbols) and at the end of the set operation (open symbols) for the twelve  $<|_{trar}, V_g>$  pairs.

is too small the RRAM cells can not perform the set transition. However, during the set operation this limitation was found in a smaller group of pairs. The overlapping of both CDFs in this situation was already clarified in the forming section. Once again, if the  $V_g$  value is too large the dispersion of the CDF becomes uncontrollable to fit different LRS levels within the read-out current range (for instance  $<20, 1.6> <\mu A, V>$ ). It is worth to point out that if the CDF shape is narrow and too close to the  $I_{trg}$  value after the set transition, as depicted in Fig. 9 for <30, 1.2> and <40, 1.4>  $<\mu$ A, V>, some cells could suffer from post-algorithm instability. Additionally, as illustrated in Fig. 10, the voltages required to perform the set operation by using these pairs are strongly increased. Therefore, a trade-off should be exercised. Extremely narrow current CDFs lead to more effective definition of discrete LRS levels within the readout current range. However, the energy consumption of the memory increases significantly because of the larger number of M-ISPVA pulses and voltages required to perform set operations, as well as, because of the error detection and correction systems required to handle the minor post-algorithm instabilities.

The evolution of the read-out currents measured during the whole M-ISPVA on a group of 128 RRAM cells are shown in Fig. 11. The results confirm the enhanced suppression of the post-algorithm instabilities compared to those observed during the forming operation. The broader range of current values covered by the RRAM devices during the M-ISPVA after the set transition by using  $\langle 30, 1.6 \rangle \langle \mu A, \rangle$ V> instead of <30, 1.4>  $<\mu A$ , V> is a further evidence of the detrimental increase of the CDF dispersion observed in Fig. 9 when the  $V_g$  is too large. The current evolution for the pair <30,  $1.2> <\mu A$ , V> shown in Fig. 11 illustrates that the Itrg value is achieved progressively after the initial current jump. This behavior explains the extremely narrow CDF shown in Fig. 9 and why several devices cross down the Itrg limit, which provoke the post-algorithm instabilities.



FIGURE 10. CDFs of the set voltages corresponding to three  $<I_{trg}$ ,  $V_g>$  pairs with  $I_{trg}$  equal to 30  $\mu$ A.



**FIGURE 11.** Evolution of the read-out current values during the M-ISPVA for the set operation on 128 RRAM devices (in grey) as a function of V<sub>g</sub> by considering  $I_{trg}$  equal to 30  $\mu$ A. The figure highlights the evolution on a specific cell.

# C. ENDURANCE

After proving the effectiveness of the M-ISPVA algorithm in countering the post-algorithm instabilities we have evaluated the multi-level endurance capabilities of the RRAM array. To this extent, we have performed 1k switching cycles (i.e., consecutive set/reset operations) by monitoring the read-out current of the four different resistance states at every cycle. This procedure is comparable to what has been used in [24]. As shown in Fig. 12, the M-ISPVA approach is not threatening the reliability of the memory since the LRSs and HRS are not degraded. Further, the post-algorithm instability is not a concern since the read window margin between different states is kept intact at the end of the programming algorithm, thus strongly reducing the number of cells moving to a wrong resistance level.

# D. LOW AND HIGH TEMPERATURE DATA RETENTION

The data retention still remains a key issue in RRAM technology [17], [21], [32]. Thermal stability in multi-level switching is even more crucial to prevent the overlapping



FIGURE 12. Endurance characterization of the four resistance levels achieved with M-ISPVA (LRSs) and ISPVA (HRS) along 1k switching cycles. For LRSs data are shown after the set transition (solid symbols) and at the end of the algorithm (open symbols). The error bars refer to the device-to-device dispersion for 128 RRAM devices.



FIGURE 13. Variation of the average and dispersion of the current values sampled during the retention test compared to the beginning of the bake at 85 °C (solid symbols) and 200 °C (open symbols) for the three LRSs conductive levels.

between adjacent LRS levels. In order to assess the low and high temperature data retention of the four conductive levels, the RRAM devices in the array were baked at 85 °C and 200 °C for 100 hours. According to one of our prior publications [17], HRS remains stable during the retention tests. Therefore, we monitored the three LRSs at different sampling times. In Fig. 13 the average and the dispersion of the read-out currents measured at each sampling time are depicted as the variation from their initial values at 0 hours ( $\Delta$ Read-out). At 85 °C the behavior is stable and only a slight decrease of the mean values occurs. In contrast, at 200 °C the thermal stress has a strong impact. The mean values feature a monotone decrease with a faster rate strongly depending on the LRS conductivity [33]. In addition, the dispersion values grow with the decrease of the LRS conductivity, which is caused by the growth of the tails of the current CDFs, as already observed in [17].

97 °C

10 years

32



FIGURE 14. Distributions of the retention failure times for the three LRSs: LRS1 (blue), LRS2 (green), and LRS3 (red); at the three annealing temperatures: 200, 230, and 260 °C.



FIGURE 15. Weibull plots of the retention failure times distributions for the three LRSs: LRS1 (blue), LRS2 (green), and LRS3 (red); at the three annealing temperatures: 200, 230, and 260 °C.

The stability of the multi-level information storage must be guaranteed for at least 10 years [34], [35]. To assess the maximum working temperature that allows such a data retention lifetime, we exploited a high temperature accelerated test [35], [36]. Sufficient statistical information from the LRSs retention failures can be obtained in a practical period of time, namely 10 hours, by using bake temperatures of 200 °C, 230 °C, and 260 °C [36], as shown in Fig. 14. As depicted in Fig. 15, the distributions of the retention failure times can be modeled by using a Weibull distribution. The retention mean time to failure (MTTF) of each LRS level programmed on the RRAM devices can be calculated for the three temperatures according to the following equation [36]:

$$MTTF = \alpha \Gamma (1 + 1/\beta) \tag{1}$$

where  $\Gamma(x)$  is the Euler's gamma function and  $\alpha$  and  $\beta$ the scale and shape parameters of the Weibull distribution, respectively. Assuming an Arrhenius dependency of the MTTF with the temperature we can state [36]:

$$MTTF = A \exp\left(\frac{E_a}{kT}\right) \tag{2}$$





three LRSs: LRS1 (blue), LRS2 (green), and LRS3 (red). The maximum temperature value that guarantees such a retention lifetime for each LRS is depicted.

where A is a pre-exponential constant, k is the Boltzmann's constant, T is the operating temperature of the device and  $E_a$  is the extracted activation energy. The three  $E_a$  values obtained for LRS1, LRS2, and LRS3 according to the data in Fig. 15 are: 1.05 eV, 1.16 eV, and 0.97 eV, respectively. These results are in line with the values already reported in literature [36] and evidence a universal mechanism associated to the physical processes involved in the degradation of the CF regardless the morphology of the CF in the conductivity range of 20-50  $\mu A$  [33], [34]. An average  $E_a$  value of 1.06 eV can be assumed to extrapolate the maximum temperature value that guarantees a MTTF of 10 years (see Fig. 16). According to Ielmini et al. [33], a stronger CF leads to a more stable LRS. However, the improvement achieved on the data retention by increasing the CF conductivity from 20 to 40  $\mu A$  is only about 5 °C, which is extremely small. In terms of multi-level operation, the temperature limit must be established as the most restrictive value, namely 97 °C.

#### **IV. CONCLUSION**

12

11

10

n(MTTF)

LRS1

LRS2 100 °C

In this study, we evaluated the feasibility of reliable multi-level operation in Al-doped HfO2 RRAM arrays by introducing the M-ISPVA algorithm. This new approach considers two parameters to shape the resistance levels:  $I_{trg}$  and Vg. Four conductive levels were successfully established in each 1T1R device of the memory. The importance of defining suitable  $\langle I_{trg}, V_g \rangle$  pairs in order to reduce post-algorithm instabilities while accomplishing endurable multi-bit operation was evidenced. The endurance up to 1k switching cycles has been demonstrated. Finally, the low and high temperature data retention were assessed reporting an  $E_a$  of 1.06 eV in accelerated tests. This value materializes in a 10 years lifetime with working temperatures below 100 °C, ensuring reliable multi-bit data storage for many NVM applications.

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